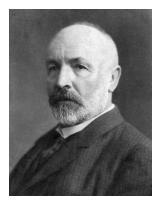
Representations of Ordinal Numbers

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September 19, 2019

Cantor

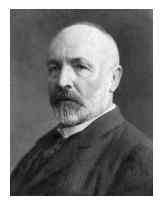


Cantor at early 20th century.*

Cantor defined ordinal numbers by two principles of generation and a first ordinal [Tiles 2004].

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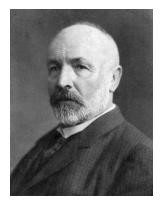
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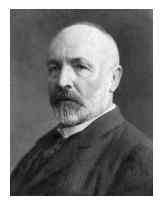
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- 0 is the first ordinal number.
- The successor of an ordinal number is an ordinal number.
- The limit of an infinite increasing sequence of ordinals is an ordinal number.

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Constructing Some Ordinals

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$$\omega^{2}, \omega^{2} + 1, \omega^{2} + 2, \dots, \omega^{3}, \omega^{3} + 1, \dots, \omega^{\omega}, \dots$$

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von Neumann Ordinals

von Neumann [1928] defined ordinals by:

Definition

An ordinal is a set α that satisfies:

- For every y ∈ x ∈ α it occurs that y ∈ α. This is called a transitive property.
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Remark

Observe that the definition is not recursive as Cantor's.

Some von Neumann Ordinals

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It is important to see that it occurs that:

$$0 \in 1 \in 2 \in \ldots \omega \in \omega + 1 \in \ldots$$

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It is important to notice that the countable ordinals are the ordinals of the first and second class of Cantor.

Hilbert Definition

Hilbert defined the natural and ordinal numbers using predicate logic [Hilbert 1925].

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$$\operatorname{Nat}(0)$$

 $\operatorname{Nat}(n) \to \operatorname{Nat}(\operatorname{succ}(n))$
 $\{P(0) \land \forall n[P(n) \to P(\operatorname{succ}(n))]\} \to [\operatorname{Nat}(n) \to P(n)]$

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$$\begin{array}{l} \operatorname{On}(0) \\ \operatorname{On}(n) \to \operatorname{On}(\operatorname{succ}(n)) \\ \{\forall n[\operatorname{Nat}(n) \to \operatorname{On}(f(n))]\} \to \operatorname{On}(\operatorname{lim}(f(n))) \\ \{P(0) \land \forall n[P(n) \to P(\operatorname{succ}(n))] \land \forall f \forall n[P(f(n)) \to P(\operatorname{lim} f)]]\} \\ \to [\operatorname{On}(n) \to P(n)] \end{array}$$

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\end{array}$$

where $Nat\ \mbox{and}\ On$ are propositional functions representing both numbers.

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The computable ordinals are less than the countable ones, as there are less λ -terms than real numbers.

The first countable ordinal that is non-computable is called ω_1^{CK*} . Furthermore, all non-countable ordinals are non-computable.

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- **0**, 1, 2, ... \rightarrow 0
- 1, **2**, 3, ... \rightarrow 1
- 2, 3, 4, ... \rightarrow 2

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- 1, **2**, 3, ... \rightarrow 1
- 2, 3, 4, ... → 2
 ...

• **0**, 2, 4, 6 ...
$$\rightarrow \omega$$

- 2, 4, 6, 8 ... $ightarrow \omega + 1$
- 4, 6, **8**, 10 ... $\rightarrow \omega + 2$

Hardy

- **0**, 1, 2, ... \rightarrow 0 • 1, **2**, 3, ... \rightarrow 1 • 2. 3. **4**. ... \rightarrow 2 • 0, 2, 4, 6 ... $\rightarrow \omega$ • 2. **4**. 6. 8 ... $\rightarrow \omega + 1$ • 4. 6. 8. 10 ... $\rightarrow \omega + 2$ • 0, 4, 8, 12, ... $\rightarrow \omega \cdot 2$ • 4, 8, 12, 16, ... $\rightarrow \omega \cdot 2 + 1$
- 8, 12, 16, 20, ... $\rightarrow \omega \cdot 2 + 2$

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$(\omega \cdot n + k)_x := 2^n (x + k)$

Martin-Löf's Representation

Martin-Löf's represented ordinals in his type theory [Martin-Löf 1984].

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	<i>n</i> : On	$f:\mathrm{Nat}\to\mathrm{On}$
zero _o : On	$succ_o n$: On	$\lim f$: On

Martin-Löf's Representation

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Martin-Löf's definition is analogous to Cantor and Hilbert's definition.

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Which ordinal cannot be constructed by Martin-Löf's representation?

Is it possible to define, similarly, a ω_1^{ML} ?

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the Theory of Ordinal Numbers". In: Fundamenta
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